

FINDING SHORTEST PATHS

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SYNOPSIS

- ◉ Purpose of Speech
- ◉ Brief Intro to Shortest Paths
- ◉ A Simple implementation: BFS
- ◉ SPFA: Optimized BFS
- ◉ Shortest Paths on Tree

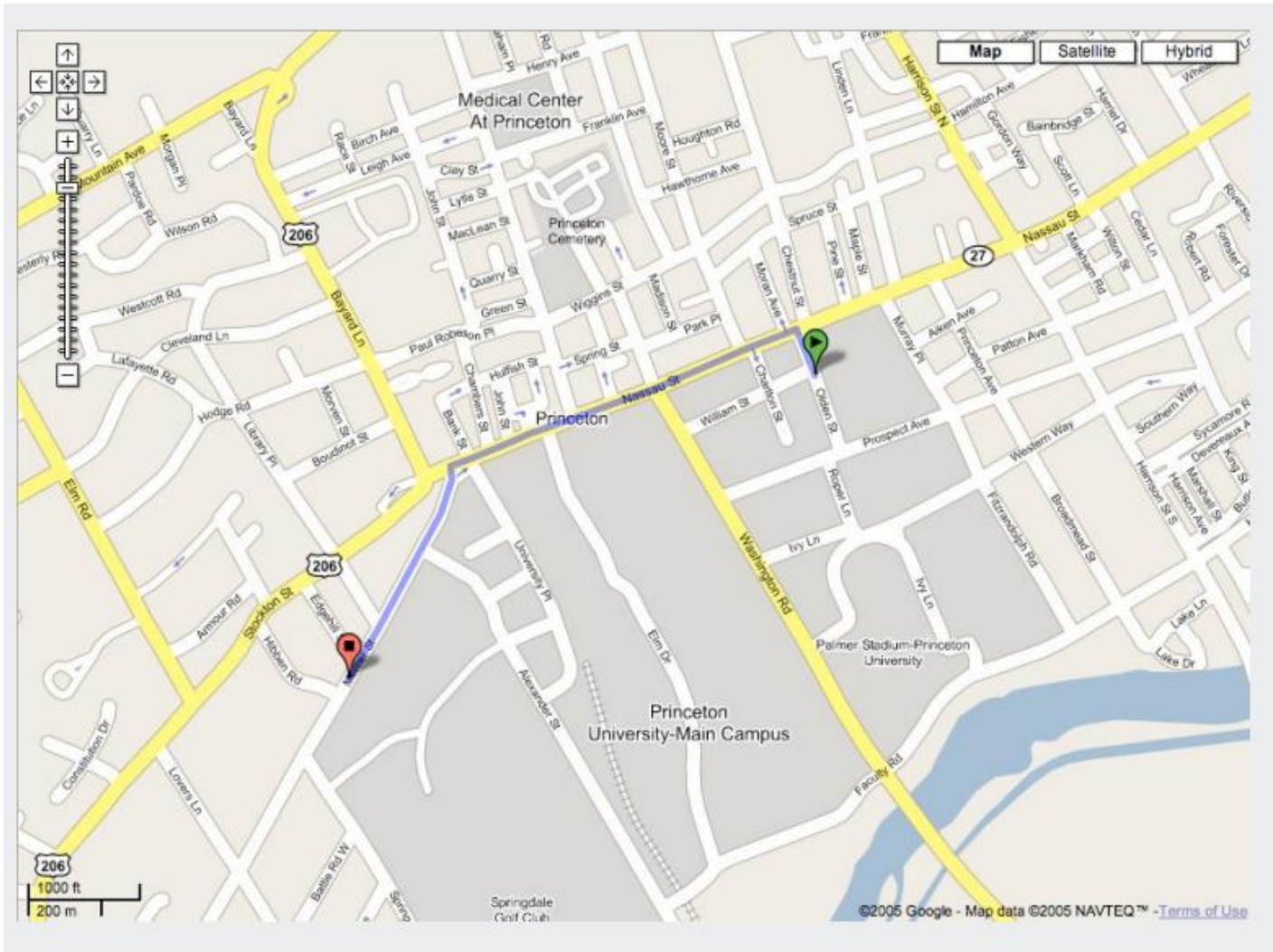
PURPOSE

- ◉ **Fundamental skills for CS students**
- ◉ **A broadly useful problem solving model**
- ◉ **Frequent appearance in Computer-based Exams**
- ◉ **OJ 2101**
- ◉ **OJ 2106**

APPLICATIONS

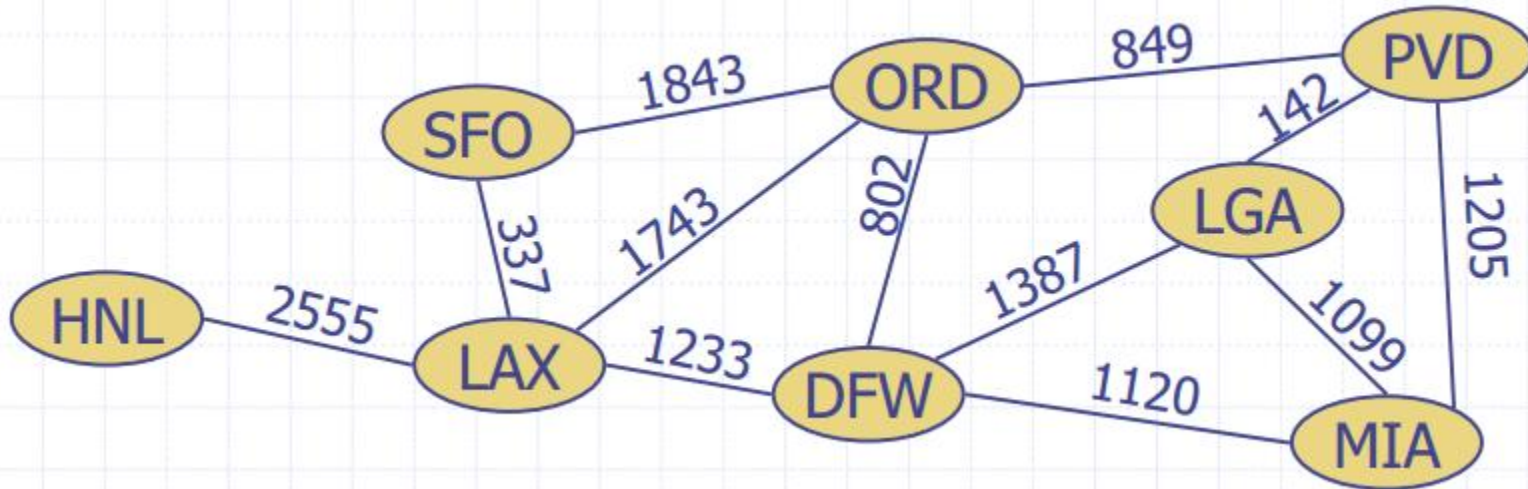
- ◉ **Maps**
- ◉ **Robot navigation**
- ◉ **Texture mapping**
- ◉ **Typesetting in Tex**
- ◉ **Urban traffic planning**
- ◉ **Optimal pipeling of VLSI chip**
- ◉ **Telemarketer operator scheduling**
- ◉ **Network routing protocols**

INTRO TO SHORTEST PATHS



INTRO TO SHORTEST PATHS

- ◉ Weighted Graphs
- ◉ Vertex
- ◉ Edge
- ◉ Distance

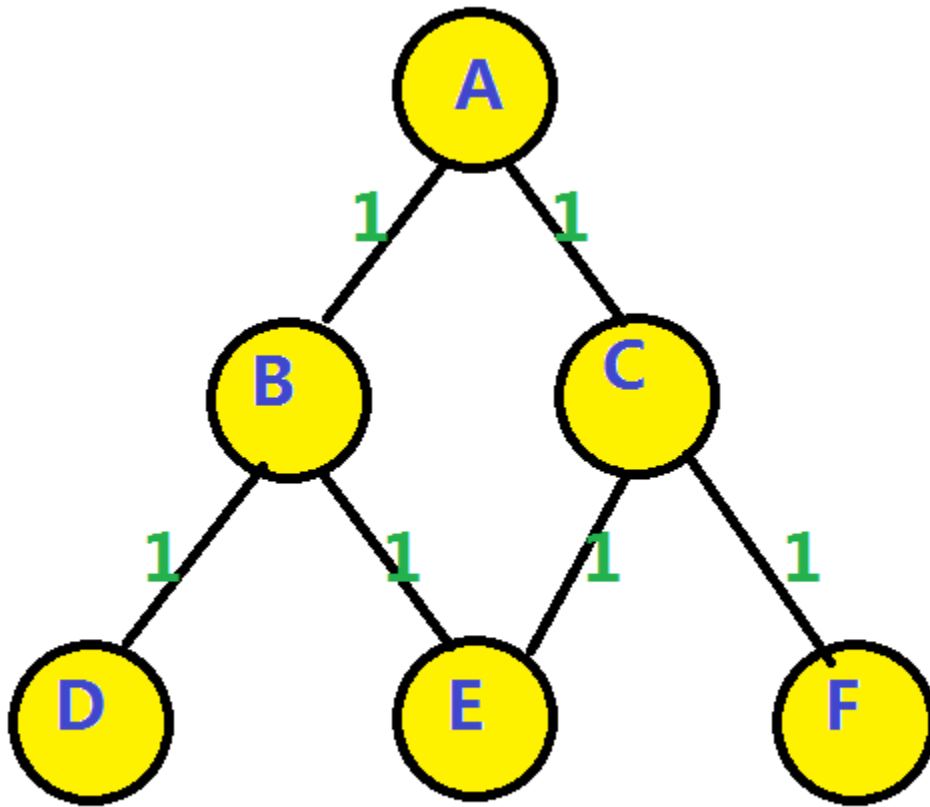


A SIMPLE IMPLEMENTATION

- ⦿ **Problem simplification:**
- ⦿ **Graph G with $|V|$ vertices and $|E|$ edges**
- ⦿ **Each edge $e=(v_1,v_2)$ has weight 1**

- ⦿ **BFS**
- ⦿ **The time when you visit the vertex must be the distance to the source.**
- ⦿ **$D[i]$**

A SIMPLE IMPLEMENTATION



Visit order(BFS):
A, B, C, D, E, F

A SIMPLE IMPLEMENTATION

- ◉ The general form:
- ◉ Graph G with $|V|$ vertices and $|E|$ edges
- ◉ Each edge $e=(v_i, v_j)$ has weight $w(i, j) > 0$

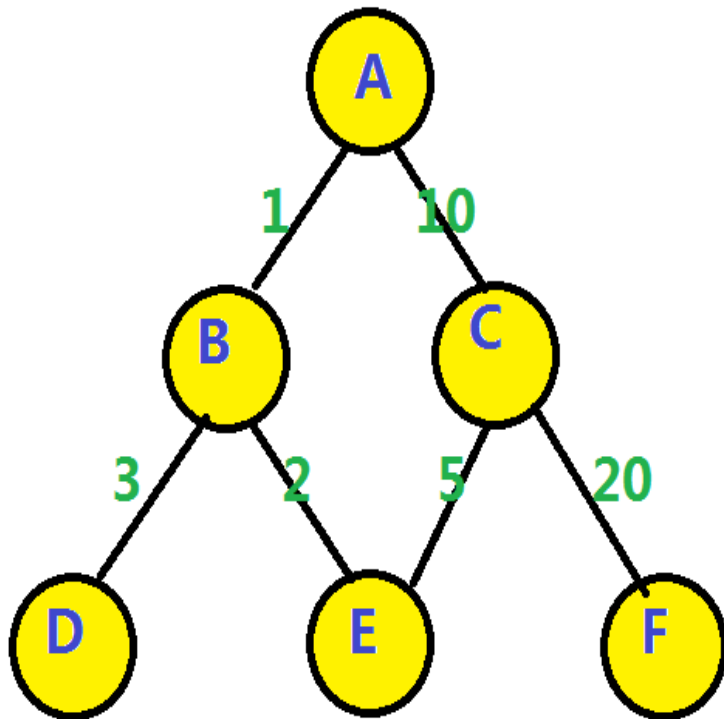
- ◉ How to compute $D[i]$
- ◉ Edge relaxation mechanism

EDGE RELAXATION

- ◉ Consider 2 vertices src and dst
- ◉ Edge $e(\text{src}, \text{dst})$ has weight $w(\text{src}, \text{dst})$

- ◉ When to relax dst?
- ◉ $D[\text{src}] + w(\text{src}, \text{dst}) < D[\text{dst}]$
- ◉ Implement within BFS

A SIMPLE IMPLEMENTATION



Visit order:

$$A:D[A] = 0$$

$$B:D[B] = 1$$

$$C:D[C] = 10$$

$$D:D[D] = 1 + 3$$

$$E:D[E] = 1 + 2$$

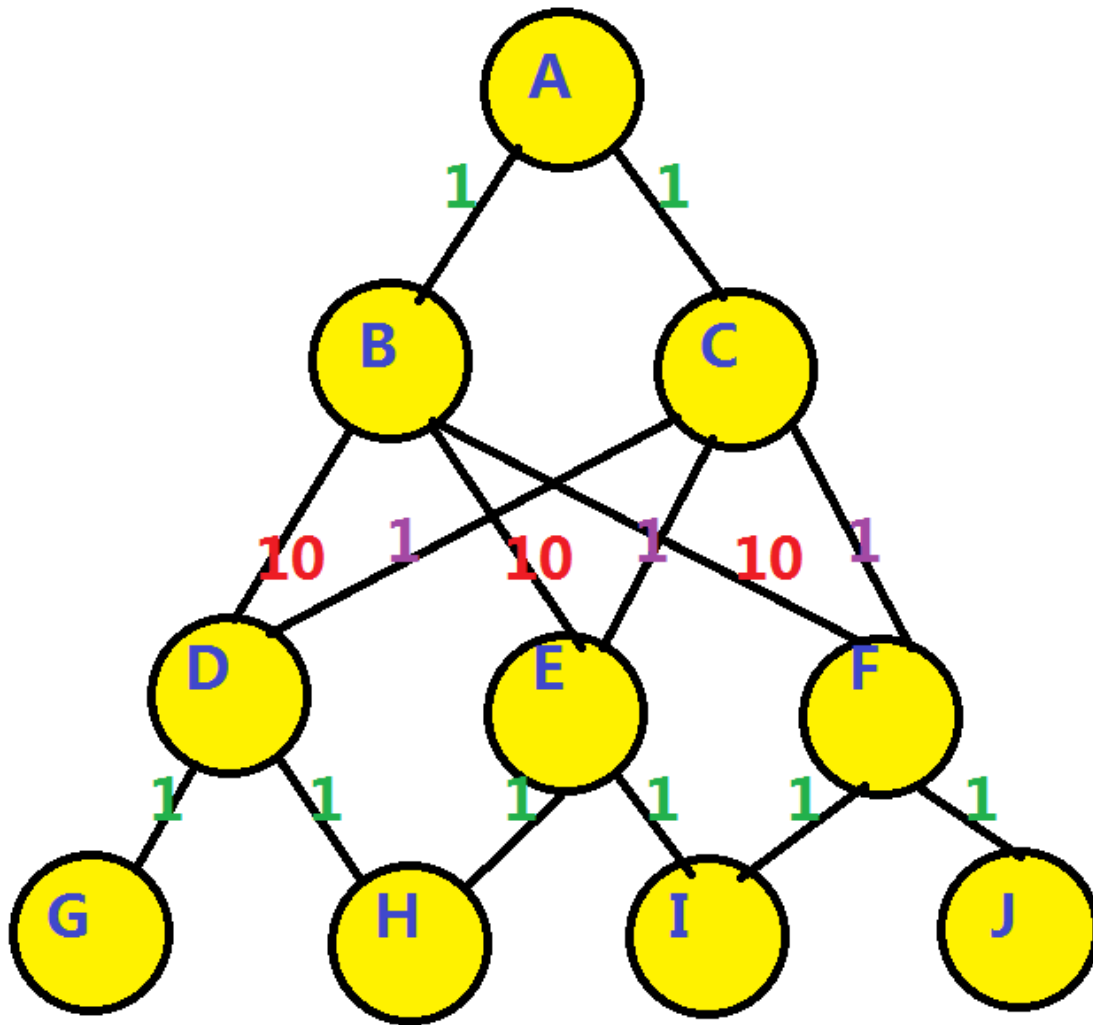
$$F:D[F] = 10 + 20$$

$$C:D[C] = 1 + 2 + 5$$

$$F:D[F] = 1 + 2 + 5 + 20$$

SPFA: OPTIMIZED BFS

- ⦿ However, you may be disappointed with the time efficiency of BFS
- ⦿ Look at the following example



⊙ A,B,C,D,E,F,D,E,F,G,H,H,I,I,J,G,H,H,I,I,J

SPFA: OPTIMIZED BFS

- ⊙ One more parameters for each point
- ⊙ $H(v)$
- ⊙ $H(v) = 1$ iff v is in the Queue
- ⊙ The Algorithm:

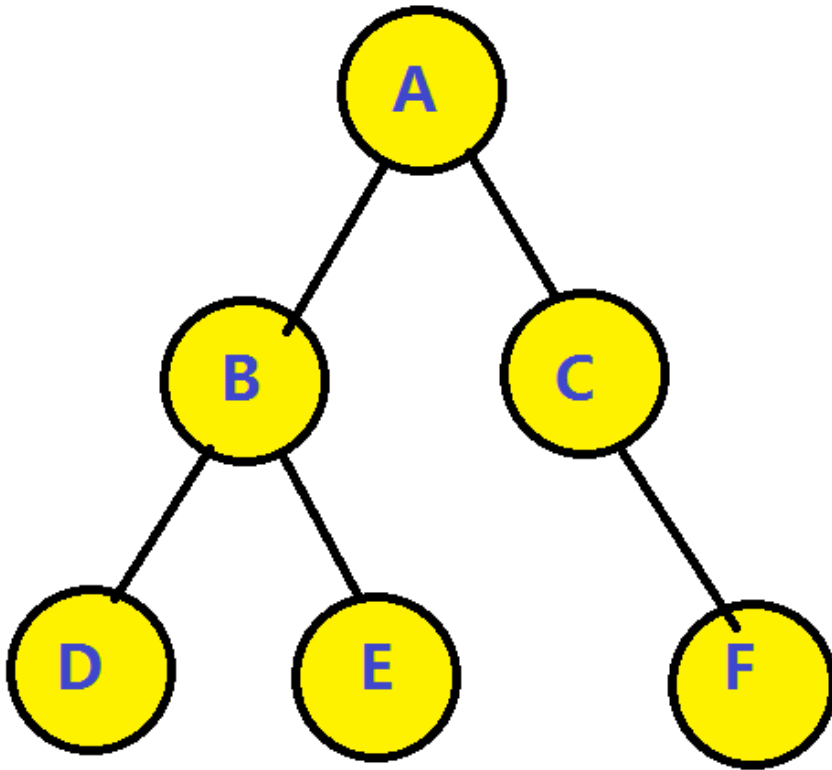
If $D[src] + w(src, dst) < D[dst]$ Then
Relax(dst)

If $H[dst] = 0$ Then `Queue.add(dst)`

Else do nothing

SHORTEST PATHS ON TREE

- ◉ **Characteristic of Tree?**



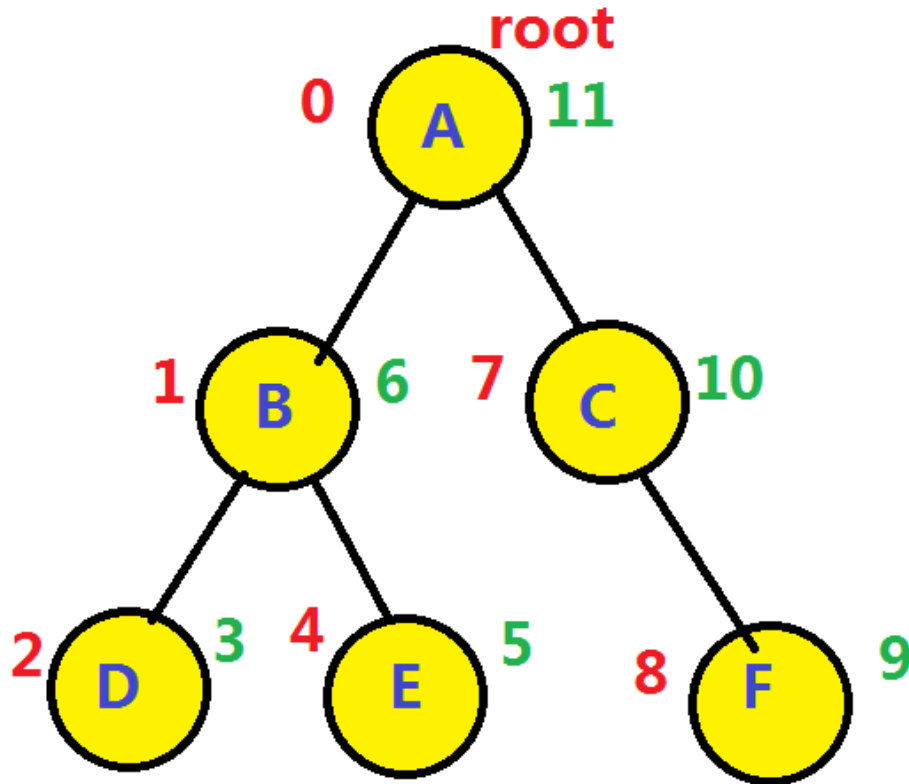
SHORTEST PATHS ON TREE

- ◉ Question?
- ◉ How to Find distance between 2 vertices on Tree?
- ◉ Time efficiency?

- ◉ Open Problem
- ◉ DFS-based implementation
- ◉ Given a rooted Tree, how to judge whether x is ancestor of y ?

SHORTEST PATHS ON TREE

◉ In[v] & Out[v]



SHORTEST PATHS ON TREE

- ◉ **DFS-based Implementation:**

- ◉ **Consider 2 vertices X and Y**

If Ancestor(X) = Y Then

$$\text{Dis}(X, Y) = \text{Dis}(\text{father}(X), Y) + w(\text{father}(X), X)$$

Else

$$\text{Dis}(X, Y) = \text{Dis}(X, \text{father}(Y)) + w(\text{father}(Y), Y)$$

- ◉ **Caution:**

If $X = Y$ Then $\text{Dis}(X, X) = 0$

THANK YOU FOR RESPECT

REFERENCES

- ◉ Princeton Online Course
- ◉ Algorithm Design